

LECTURE 7

MONDAY JANUARY 27

transitions

$q_0 : \$ > q_1$

$\boxed{f(q_0, q_1)} : \$ > q_1 \quad X$

(Compile time)

Concrete syntax 1

website URL
states
"

Abstract syntax

Runtime
(Intermediate format)

fa

Epsilon NFA
DFA

Subject
Conj.

fa

DFA

Concrete syntax 2

machine M

States =

$\{q_0, q_1, q_f\}$

Transitions =

$\{(q_0, e, \{q_1, q_f\})\}$

!

From RE to Scanner (1)

Regular Expression: $r[0..9]^*$

Token Type (CharCat)

r	0, 1, 2, ..., 9	EOF	Other
Register	Digit	Other	Other

Transition

Register	Digit	Other
s_0	s_1	s_e
s_1	s_e	s_2
s_2	s_e	s_e
s_e	s_e	s_e



Token Type (Type)

s_0	s_1	s_2	s_e
invalid	invalid	register	invalid

NextWord()

```
-- Stage 1: Initialization  
state :=  $s_0$ ; word :=  $\epsilon$   
initialize an empty stack s; s.push(bad)  
-- Stage 2: Scanning Loop  
while state ≠  $s_e$   
    NextChar(char); word := word + char  
    if state ∈ F then reset stack s end  
    s.push(state)  
    cat := CharCat(char)  
    state :=  $\delta$ [state, cat]  
-- Stage 3: Rollback Loop  
while state ∈ F ∧ state ≠ bad  
    state := s.pop()  
    truncate word  
-- Stage 4: Interpret and Report  
if state ∈ F then return Type[state]  
else return invalid  
end
```

Example input: r241

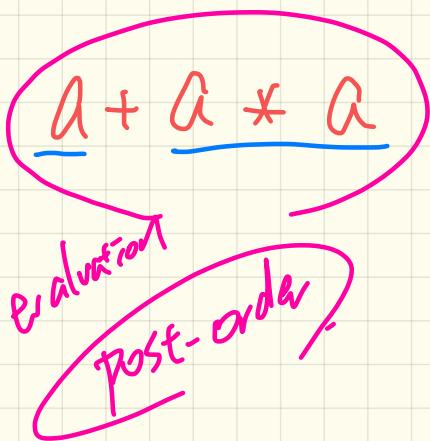
state : $s_0 \xrightarrow{r} s_1 \xrightarrow{2} s_2 \xrightarrow{4} s_1 \xrightarrow{1} s_e$

word : r241 ~~bad~~

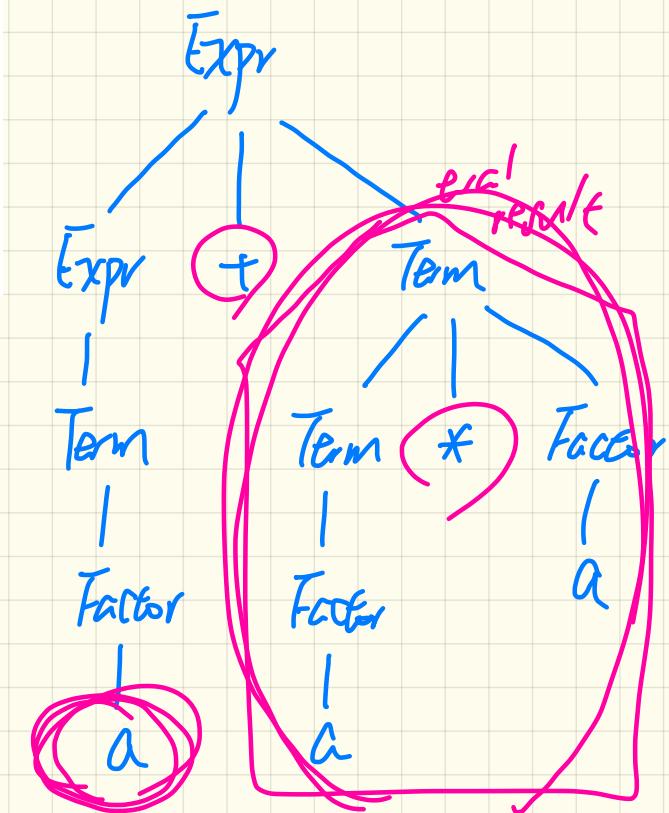
word: r241
syn. cat: register

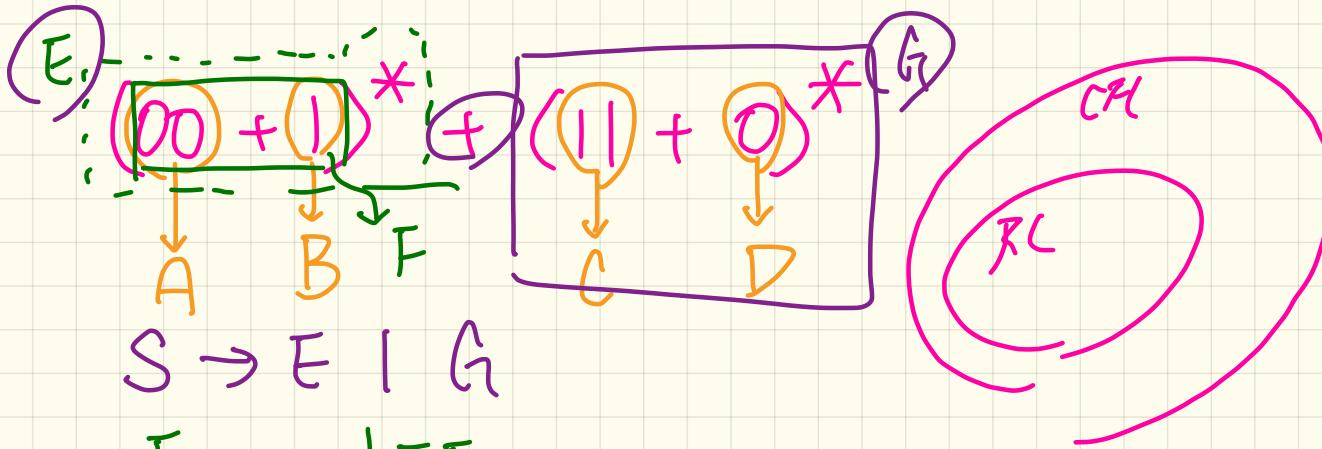
bad

$\text{Expr} \rightarrow \text{Expr} + \text{Term}$
 |
 $\text{Term} \rightarrow \text{Term} * \text{Factor}$
 |
 $\text{Factor} \rightarrow (\text{Expr})$
 |
 a



 Evaluation post-order





$S \rightarrow E \mid G$

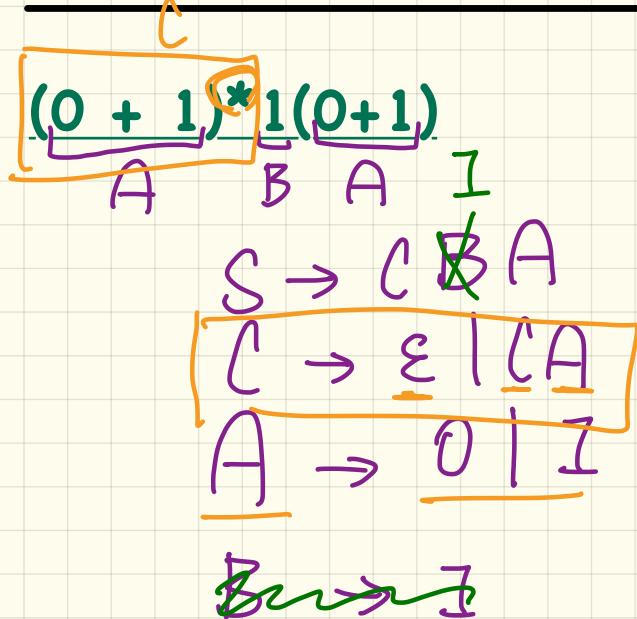
$E \rightarrow \epsilon \mid FE$
 $F \rightarrow A \mid B$
 $A \rightarrow 00$

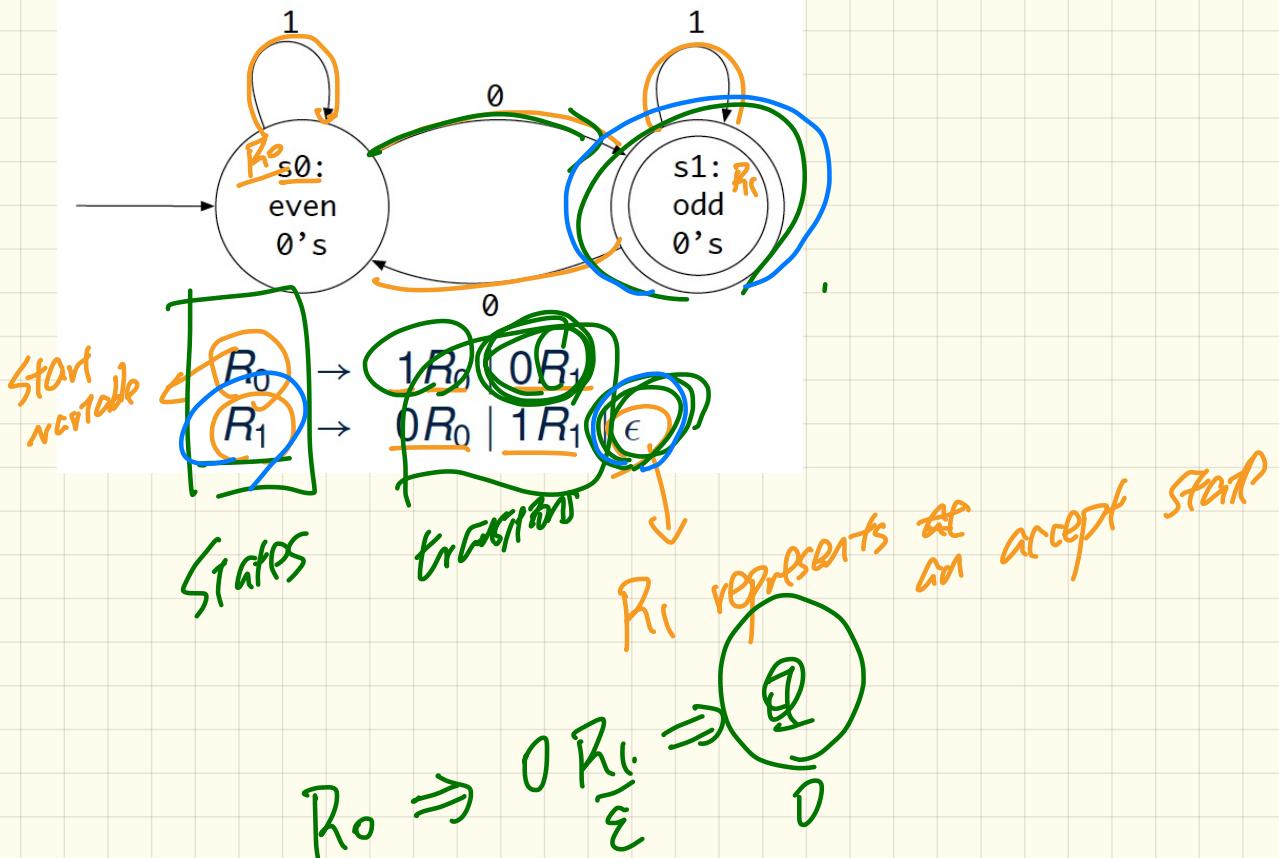
$B \rightarrow 1$

$C \rightarrow 11$

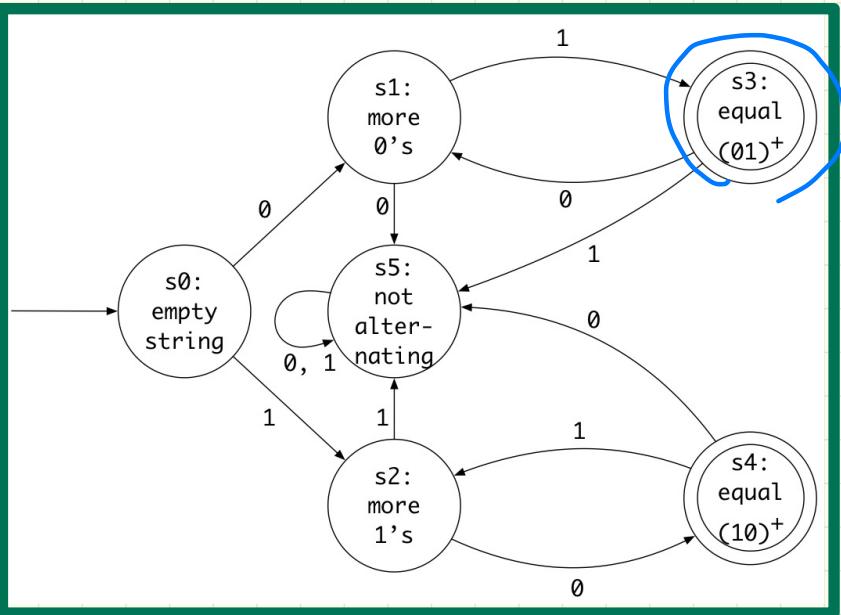
$D \rightarrow 0$

Context-Free Grammar (CFG): from RE





Context-Free Grammar (CFG): from DFA



Start var. ab
S₀ →

S₁ →

S₂ →

S₃ → $\epsilon | 0S_1 | 1S_2$

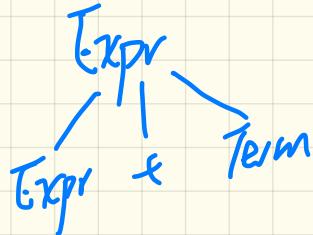
S₄ → ϵ

S₅ → $0S_1 | 1S_2$

Context-Free Grammar (CFG): Leftmost Derivation

<u>Expr</u>	\rightarrow	Expr	$+$	Term
		Term		
<u>Term</u>	\rightarrow	Term	$*$	Factor
		Factor		
<u>Factor</u>	\rightarrow	(Expr)
		a		

Parse Tree: $a + a * a$



Derivation: $a + a^* a$

$\text{Expr} \Rightarrow \underline{\text{Expr}} + \text{Term}$
 $\Rightarrow \text{Term} + \text{Term}$
 $\Rightarrow \boxed{\underline{\text{Factor}} + \text{Term}}$
 $\Rightarrow \underline{a} + \boxed{\text{Term}}$
 $\Rightarrow a + \underline{\text{Term}} * \text{Factor}$
 $\Rightarrow a + \underline{\text{Factor}} * \text{Factor}$
 $\Rightarrow a + a * \underline{\text{Factor}}$
 $\Rightarrow a + a * a$

Context-Free Grammar (CFG): Rightmost Derivation

<i>Expr</i>	\rightarrow	<i>Expr</i>	$+$	<i>Term</i>
		<i>Term</i>		
<i>Term</i>	\rightarrow	<i>Term</i>	$*$	<i>Factor</i>
		<i>Factor</i>		
<i>Factor</i>	\rightarrow	<i>(Expr)</i>		
		a		

Derivation: a + a * a

Parse Tree: a + a * a

Context-Free Grammar (CFG): Leftmost Derivation

<i>Expr</i>	\rightarrow	<i>Expr</i>	$+$	<i>Term</i>
		<i>Term</i>		
<i>Term</i>	\rightarrow	<i>Term</i>	$*$	<i>Factor</i>
		<i>Factor</i>		
<i>Factor</i>	\rightarrow	<i>(Expr)</i>		
		a		

Derivation: $(a + a)^* a$

Parse Tree: a + a * a

Context-Free Grammar (CFG): Rightmost Derivation

<i>Expr</i>	\rightarrow	<i>Expr</i>	$+$	<i>Term</i>
		<i>Term</i>		
<i>Term</i>	\rightarrow	<i>Term</i>	$*$	<i>Factor</i>
		<i>Factor</i>		
<i>Factor</i>	\rightarrow	<i>(Expr)</i>		
		a		

Derivation: $(a + a)^* a$

Parse Tree: a + a * a

Given input string $w \in \Sigma^*$

derivation
 \neq
derivation Σ

w



G is ambiguous.

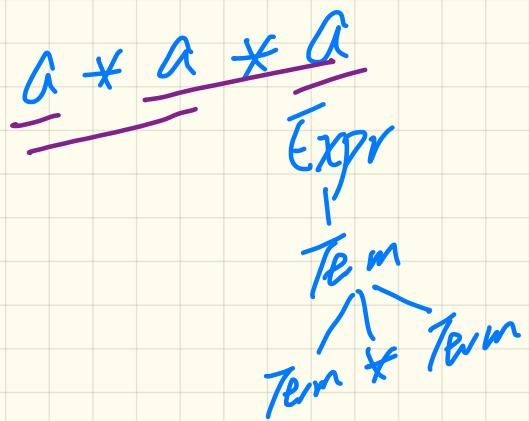
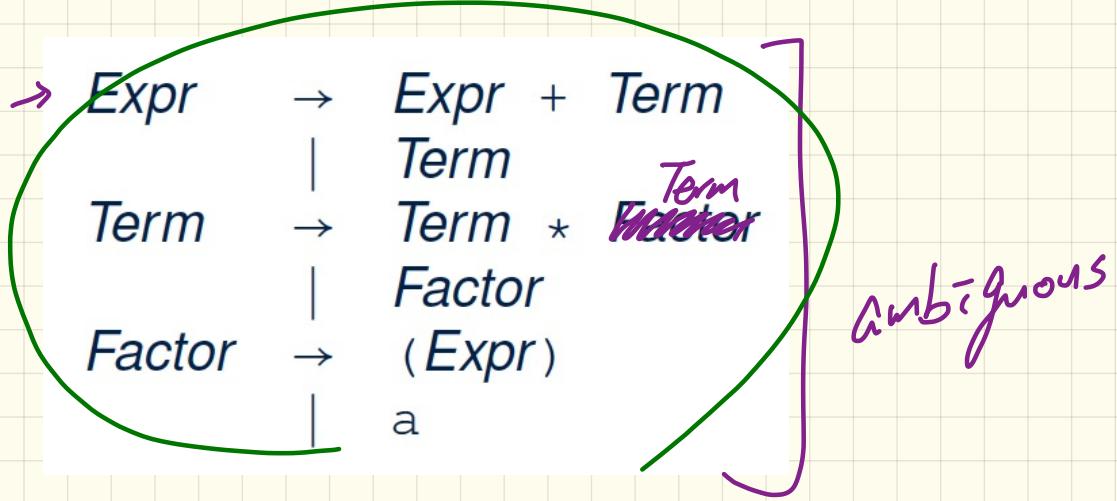
$$(\boxed{a} + \boxed{a}) * \boxed{a}$$

Exercise

ex. $a + a * a$] meaning /

ex. $a + (a * a)$] meaning 2.

ex. $(a + a) * a$] do they have same meaning according to the rule

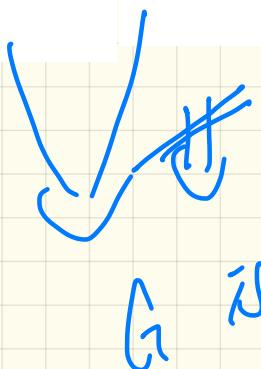


Unique leftmost derivation for the string $(a + a) * a$:

$\begin{aligned} \text{Expr} &\Rightarrow \text{Term} \\ &\Rightarrow \text{Term} * \text{Factor} \\ &\Rightarrow \text{Factor} * \text{Factor} \\ &\Rightarrow (\text{Expr}) * \text{Factor} \\ &\Rightarrow (\text{Expr} + \text{Term}) * \text{Factor} \\ &\Rightarrow (\text{Term} + \text{Term}) * \text{Factor} \\ &\Rightarrow (\text{Factor} + \text{Term}) * \text{Factor} \\ &\Rightarrow (a + \text{Term}) * \text{Factor} \\ &\Rightarrow (a + \text{Factor}) * \text{Factor} \\ &\Rightarrow (a + a) * \text{Factor} \\ &\Rightarrow (a + a) * a \end{aligned}$

Unique rightmost derivation for the string $(a + a) * a$:

$\begin{aligned} \text{Expr} &\Rightarrow \text{Term} \\ &\Rightarrow \text{Term} * \text{Factor} \\ &\Rightarrow \text{Term} * a \\ &\Rightarrow \text{Factor} * a \\ &\Rightarrow (\text{Expr}) * a \\ &\Rightarrow (\text{Expr} + \text{Term}) * a \\ &\Rightarrow (\text{Expr} + \text{Factor}) * a \\ &\Rightarrow (\text{Expr} + a) * a \\ &\Rightarrow (\text{Term} + a) * a \\ &\Rightarrow (\text{Factor} + a) * a \\ &\Rightarrow (a + a) * a \end{aligned}$



G is ambiguous
- i.e. there are two different derivations
in grammar order.

dangling

else -

